

When do I need an invariant?

CS 5010 Program Design Paradigms
“Bootcamp”
Lesson 7.4



Learning Objectives

- At the end of this lesson, the student should be able to
 - decide whether a purpose statement needs an invariant or not.

When do I need an invariant?

- It all depends on the purpose statement.
- If your code fulfills the purpose statement for any arguments of the types listed in the contract, you don't need an invariant.
- If the function fulfills its purpose statement only for certain values or combinations of values of the arguments, then you must document that restriction with a WHERE-clause.

What kind of thing belong in an invariant?

- If the function needs additional information that is not in the arguments, then you need an invariant to document the needed information
- What kind of information might you want?
 - context information (e.g. we are position n in the list)
 - other knowledge that isn't expressed in the contract (e.g. we've figured out the ball isn't going to bounce).

Whose responsibility is it?

- The invariant, along with the contract, sets down the assumptions that each function makes about the arguments that it processes
- It is up to each caller of the function to make sure that the invariant is true at every call.
- The function gets to assume that the invariant is true.

Example:

```
;; ball-normal-motion : Ball -> Ball
```

```
;; GIVEN: a Ball
```

```
;; RETURNS: the state of the ball after a
```

```
;; tick.
```

```
(define (ball-normal-motion b)
```

```
  (make-ball
```

```
    (+ (ball-x-pos b) BALLSPEED)))
```

Doesn't work for every Ball!..
Needs more information

Invariant provides the
necessary information

Example

```
;; number-list-from : ListOfX Number -> NumberedListOfX
;; RETURNS: a list with same elements as lst, but numbered
;; starting at n.
;; EXAMPLE: (number-list-from (list 88 77) 2)
;;          = (list (list 2 88) (list 3 77))
;; STRATEGY: Use template for ListOfX on lst
(define (number-list-from lst n)
  (cond
    [(empty? lst) empty]
    [else
     (cons
      (list n (first lst))
      (number-list-from (rest lst) (+ n 1)))]))
```

Works for any lst and n, so
no invariant necessary.

Function can't fulfill its purpose unless it knows where `slst` is in `lst0`

Same Code, different use statement

```
;; number-list-from :  
;;   ListOfX Number -> NumberedListOfX  
;; GIVEN: a sublist slst of some list lst0  
  
;; RETURNS: a copy of slst numbered according to its  
;; position in lst0.  
;; STRATEGY: Use template for ListOfX on slst  
(define (number-sublist slst n)  
  (cond  
    [(empty? slst) empty]  
    [else  
     (cons  
      (list n (first slst))  
      (number-sublist (rest slst) (+ n 1))))])
```

Invariant supplies the extra information

Wait, weren't those functions very similar?

- Yes. In fact they were identical (except for their names).
- The moral of the story is that it is the purpose statement that determines whether you need an invariant.

Once more: When do I need an invariant?


- If your code fulfills the purpose statement for any arguments of the types listed in the contract, you don't need an invariant.
- If the function only works for certain values or combinations of values of the arguments, then you must document the assumptions that it needs with a WHERE-clause (i.e. an invariant).

What needs to be in my purpose statement?

- The purpose statement must account for all the parameters.
 - if it doesn't then either you are passing more parameters than you need, or there's something going on that you haven't described.
- The RETURNS clause must describe the value returned by the function for all possible values of the parameters.
- If the RETURNS clause describes the value returned by the function only for some values of the arguments or some combination of arguments, then that restriction must be stated in a WHERE clause.
- It becomes the responsibility of the caller to guarantee that the restriction is satisfied.

Another example

```
;; add-remaining-length : LoN -> LoN
;; RETURNS: a list like the original, but with each
;; element increased by the length of the sublist
;; starting at that element.
;; (100 300 500) => (103 302 501)
;; Strategy: SD on lst
(define (add-remaining-length lst)
  (cond
    [(empty? lst) empty]
    [else (cons
            (+ (first lst) (length lst))
             (add-remaining-length
              (rest lst)))])))
```



Yuck!

Let's help the function along by giving it the length of the list as an argument

```
;; add-remaining-length-1 : LoN Number-> LoN
;; GIVEN: a Lon lst and a number n
;; WHERE: n = (length lst)
;; RETURNS: a list like the original, but with each
;; element increased by the length of the sublist
;; starting at that element.
;; (100 300 500) 3 => (103 302 501)
;; Strategy: SD on lst
(define (add-remaining-length-1 lst n)
  (cond [(empty? lst) empty]
        [else (cons
                (+ (first lst) n)
                (add-remaining-length-1 (rest lst)
                                       (- n 1))))]))
```

Doesn't give the right answer unless invariant is satisfied

Summary: When do I need an invariant?

- It all depends on your purpose statement!
- If the function needs additional information that is not in the arguments, then you need an invariant to document the needed information
- It is up to each caller of the function to make sure that the invariant is true at every call.

Summary

- The student should now be able to
 - decide whether a purpose statement needs an invariant or not.

Next Steps

- If you have questions about this lesson, ask them on the Discussion Board
- Go on to the next lesson